Performance Evaluation of TCP BBRv2 Alpha for Wired Broadband, considering Buffer Sizes, Packet Loss Rates, RTTs, and Number of Flows

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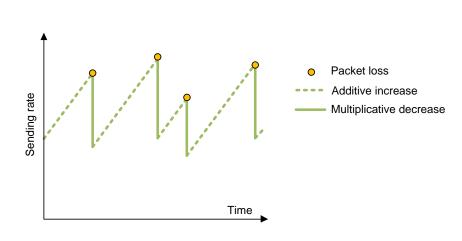
We want to recognize Neal Cardwell for his helpful comments and suggestions

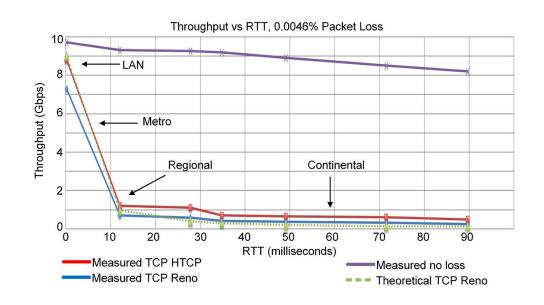
Agenda

- TCP loss-based congestion control
- Bottleneck Bandwidth and Round-trip propagation time (BBR)
- BBRv1
- BBRv2 and design goals
- Experimental setup
- Results and evaluations
- Conclusions

TCP Loss-based Congestion Control (CC)

- The principles of loss-based CC were described in the 1980s¹
- Traditional CC algorithms follow the additive-increase multiplicative-decrease (AIMD) form of congestion control





1. V. Jacobson, M. Karels, Congestion avoidance and control, ACM SIGCOMM Computer Communication Review 18 (4) (1988).

BBRv1

- BBRv1 was published in 2016
- It produces high throughput in the presence of packet losses
- However, it shows some performance issues
 - Unfairness with loss-based CC (e.g., Reno, CUBIC)
 - Low throughput for loss-based CC
 - ➤ High packet loss rates when the buffer is smaller than 1.5 bandwidth-delay product (BDP)
 - ➤ Throughput variation

^{1.} N. Cardwell et al. "BBR v2, A Model-based Congestion Control." IETF 104, March 2019.

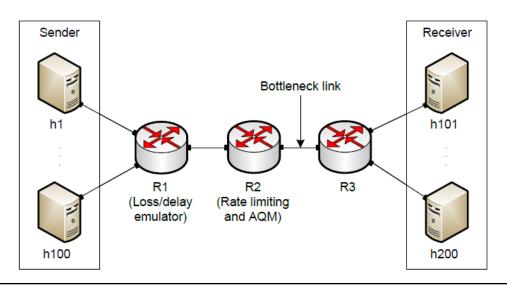
BBRv2 Design Goals

- Improving co-existence with loss-based when sharing a bottleneck link
- Avoiding the bufferbloat problem
- Minimizing the time to reach an equilibrium point (fairness)
- Reducing the variation of the throughput
- Producing high throughput, even with moderate packet losses

^{1.} N. Cardwell et al. "BBR v2, A Model-based Congestion Control." IETF 104, March 2019.

Experimental Setup

- Up to 100 simultaneous flows
- Mininet network, Linux protocol stack, Tail drop AQM policy (default)
- 300 seconds experiment duration, average results are reported
- CPU cores (8 Xeon 6130) are overprovisioned (below 50% utilization); idle CPU time
 is a good fidelity indicator, as it indicates that a host is not starved for CPU resources¹

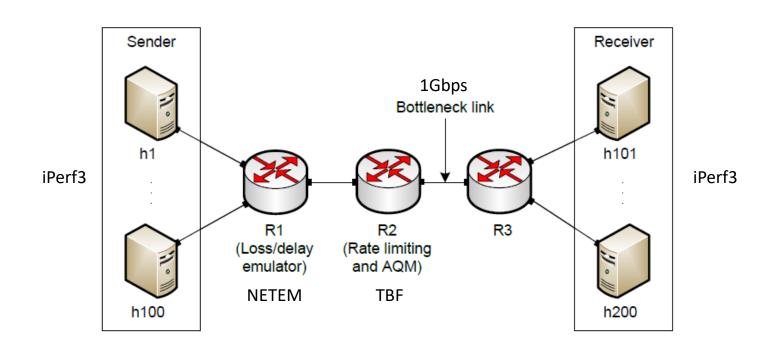


¹N. Handigol et al., "Reproducible Network Experiments Using Container-Based Emulation," ACM CoNEXT '12.



Experimental Setup

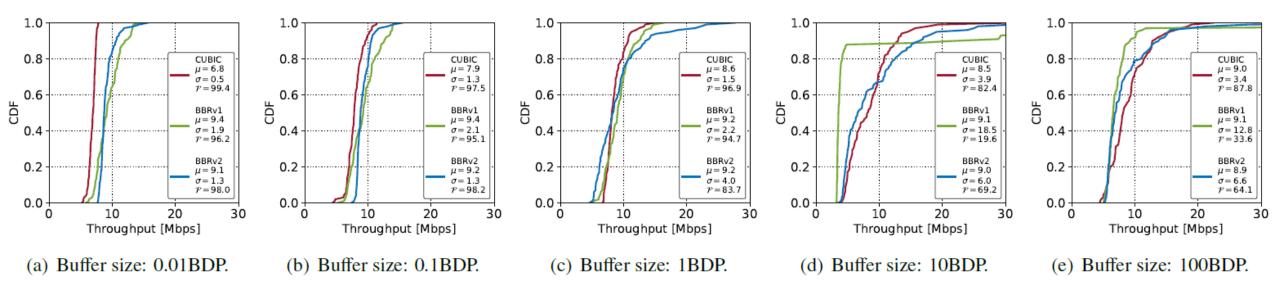
- BBRv2 version: v2alpha-2019-07-28
- Bottleneck link = 1Gbps
- All other links = 40Gbps
- Tools: iPerf3, netstat, ss, tc [tbf | netem | fq_codel | cake]



South Carolina

Multiple Flows and Buffer Sizes

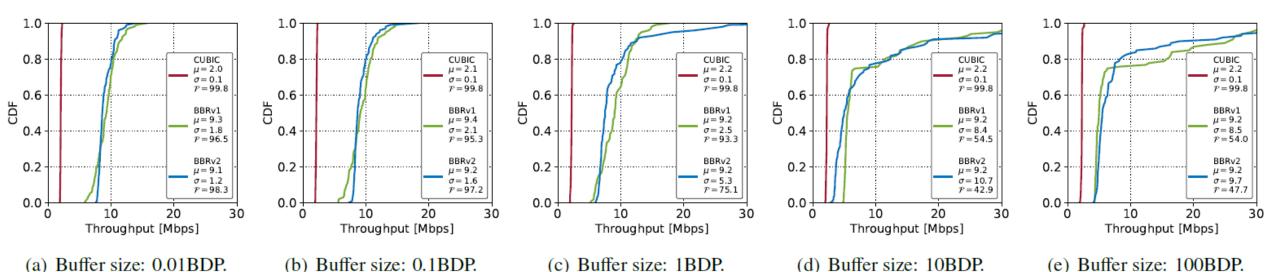
- 100 simultaneous flows, 20ms propagation delay, same CC
- Bottleneck is 1Gbps (ideal allocation is 10Mbps per flow)
- Cumulative distribution function, mean, standard deviation, and fairness



No emulated packet losses

Multiple Flows and Buffer Sizes

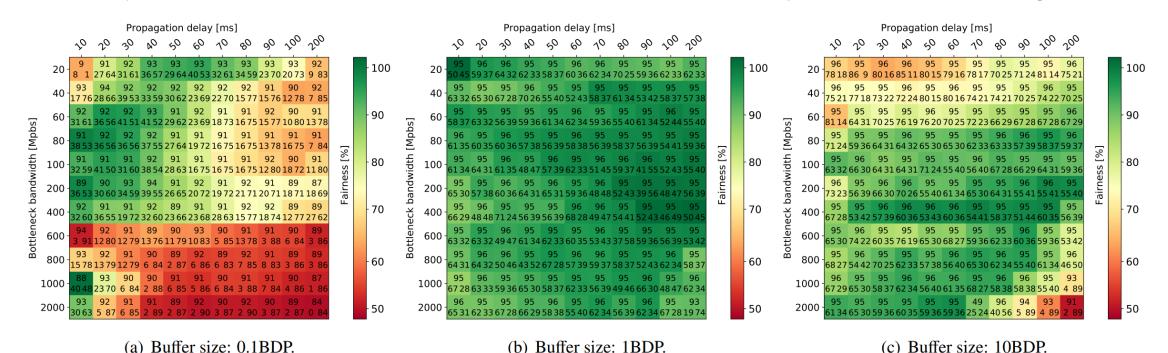
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1% packet loss rate

Coexistence and Fairness with CUBIC

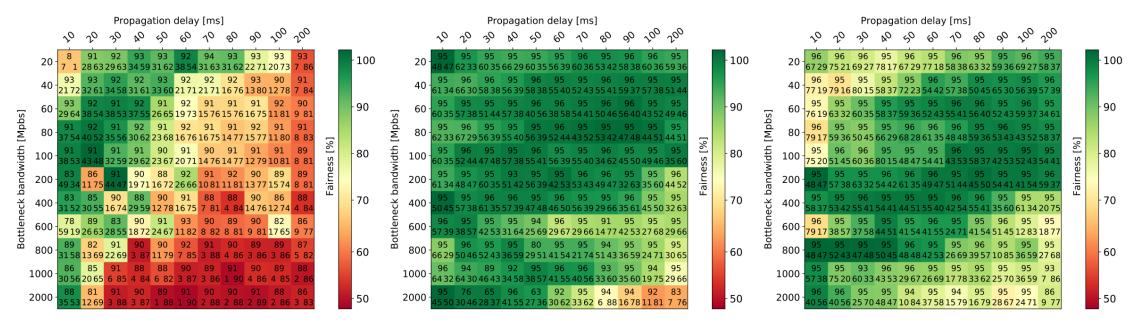
- Fairness index visualized in a heatmap for CUBIC and BBRv2 flows
- 1 CUBIC vs. 1 BBRv2
- Each entry has 3 numbers (in percentage): link utilization (center top), bandwidth used by CUBIC (bottom left value), and bandwidth used by BBRv2 (bottom right)



No emulated packet losses

Coexistence and Fairness with CUBIC

- Fairness index visualized in a heatmap for CUBIC and BBRv2 flows
- 1 CUBIC vs. 1 BBRv2
- Each entry has 3 numbers (in percentage): link utilization (center top), bandwidth used by CUBIC (bottom left value), and bandwidth used by BBRv2 (bottom right)



(d) Buffer size: 0.1BDP.

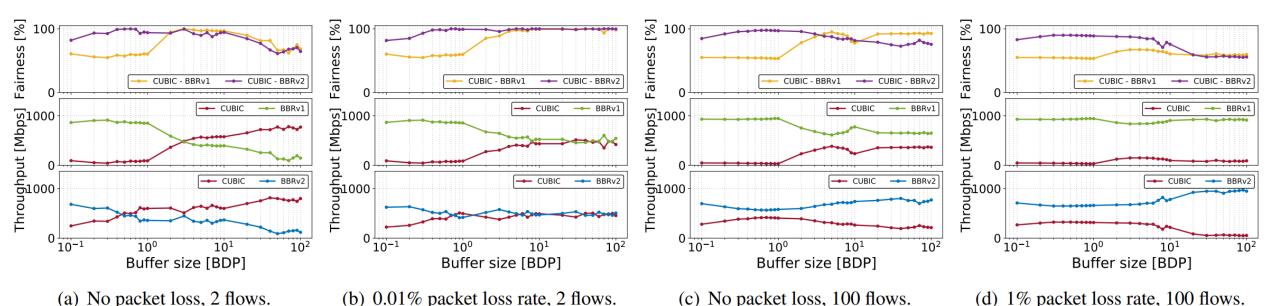
(e) Buffer size: 1BDP.0.01% packet losses

(f) Buffer size: 10BDP.



Coexistence and Fairness with CUBIC

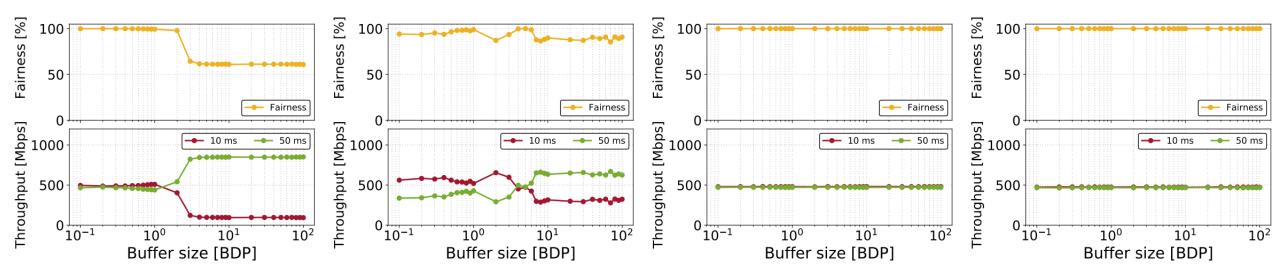
- Bottleneck is 1Gbps, 20ms propagation delay
- Competing flows sharing the same bottleneck link
- BBRv2 shows a better coexistence with CUBIC than BBRv1



Throughput and fairness as functions of the buffer size

Round-trip Time Unfairness

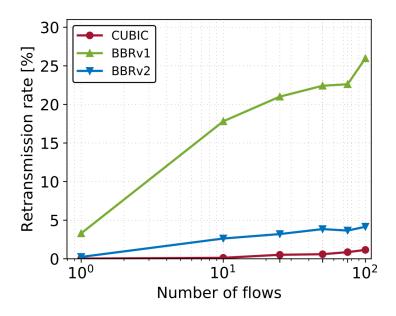
- Two simultaneous flows, one with 10ms propagation delay, and the other with 50ms
- Bottleneck is 1Gbps
- Tail Drop and FQ-CoDel AQMs
- FQ-CoDel mitigates the RTT unfairness



- (a) BBRv1 with Tail Drop policy.
- (b) BBRv2 with Tail Drop policy.
- (c) BBRv1 with FQ-CoDel policy.
- (d) BBRv2 with FQ-CoDel policy.

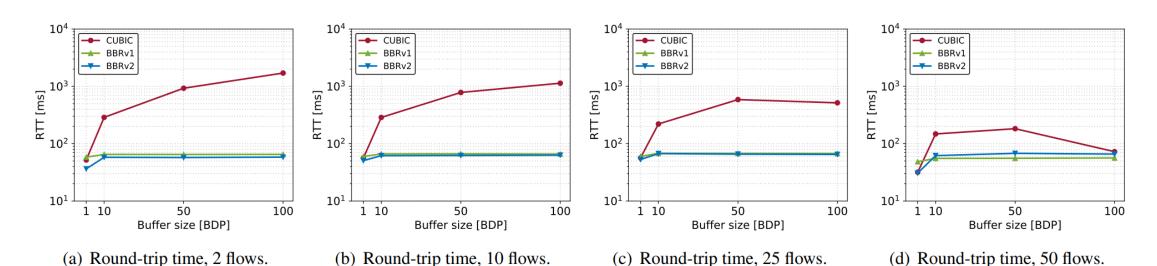
Retransmissions – Number of Flows

- The buffer size is 0.02BDP (i.e., 2.5MB, ~1,666 packets)
- Bottleneck is 1Gbps, 100ms propagation delay
- BBRv1: high retransmission rate with any number of flows (2.5%-25%)
- BBRv2: lower retransmission rate (~2.5%)



Queueing Delay with Large Buffers

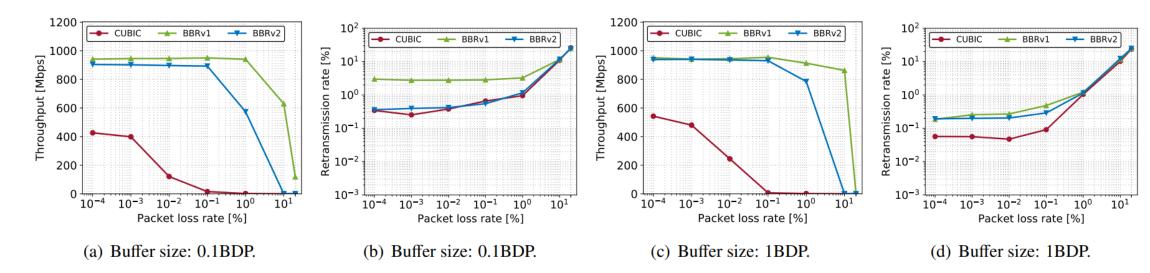
- 2, 10, 25, and 50 simultaneous flows, same CC
- Bottleneck is 1Gbps, 30ms propagation delay
- RTT: queueing delay + propagation delay
- BBRv1 and BBRv2 have low queuing delay independently of the number of flows



Round-trip time experienced with simultaneous flows

Throughput, Retransmissions – Packet Loss Rate

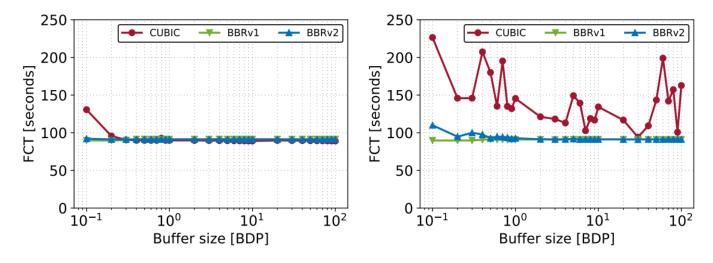
- Single flow
- Bottleneck is 1Gbps, 100ms propagation delay
- BBRv1 is loss-agnostic, which leads to a higher retransmission rate



Throughput and retransmission rate as functions of the packet loss rate

Flow-Completion Time (FCT)

- 100 simultaneous flows, same CC
- Bottleneck is 1Gbps, 20ms propagation delay
- Data transfer of 10 GB

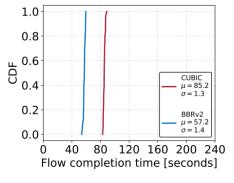


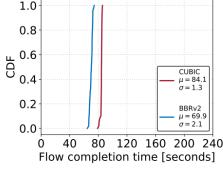
- (a) No random packet losses.
- (b) 0.01% random packet loss rate.

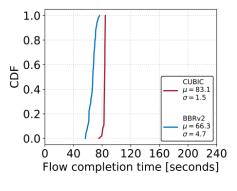
Flow completion time as a function of the buffer size

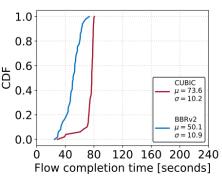
Flow-Completion Time (FCT)

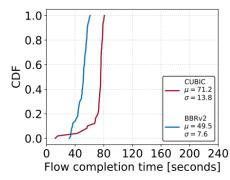
- Bottleneck is 1Gbps, 20ms propagation delay, 100 simultaneous flows
- 50 CUBIC, 50 BBRv1/BBRv2, each flow completing 100MB of data transfer

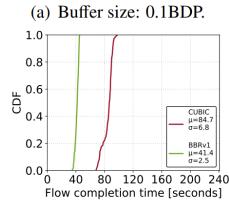




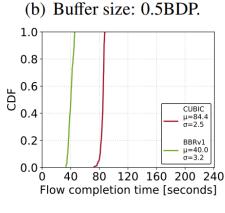


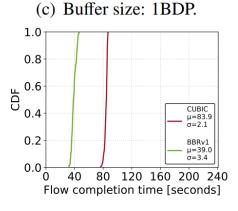


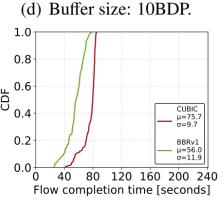


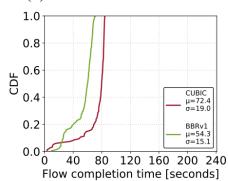


(a) Buffer size: 0.1BDP.









(e) Buffer size: 100BDP.

(b) Buffer size: 0.5BDP.

(c) Buffer size: 1BDP.

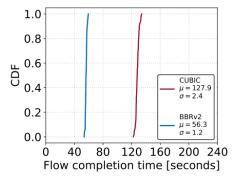
No losses

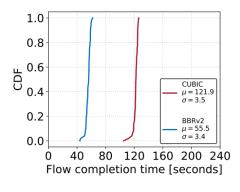
(d) Buffer size: 10BDP.

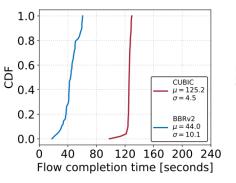
(e) Buffer size: 100BDP.

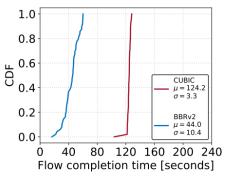
Flow-Completion Time (FCT)

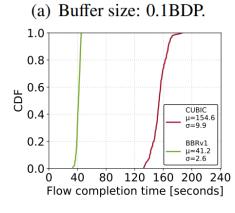
- Bottleneck is 1Gbps, 20ms propagation delay, 100 simultaneous flows
- 50 CUBIC, 50 BBRv1/BBRv2, each flow completing 100MB of data transfer



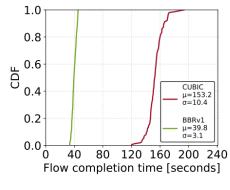




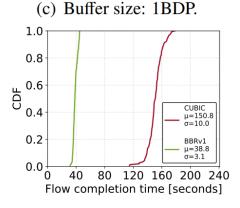


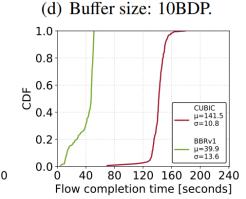


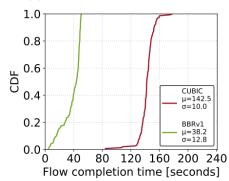
(a) Buffer size: 0.1BDP.



(b) Buffer size: 0.5BDP.







(e) Buffer size: 100BDP.

(b) Buffer size: 0.5BDP.

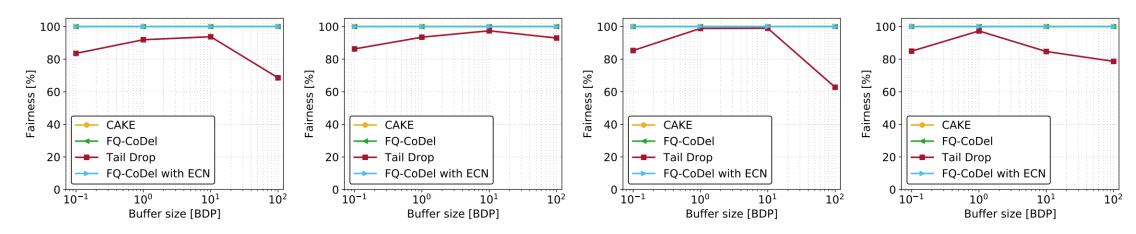
(c) Buffer size: 1BDP.1% packet losses

(d) Buffer size: 10BDP.

(e) Buffer size: 100BDP.

Impact of AQM on Fairness

- Bottleneck is 1Gbps, 20ms propagation delay
- 2, 11, 50 and 100 simultaneous flows, uneven composition
- Tail drop, CAKE, FQ_CoDel, FQ_CoDel w/ ECN

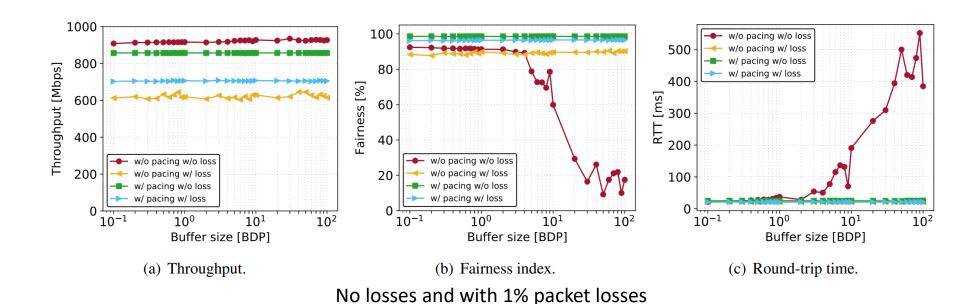


- (a) 1 BBRv2 flow, 1 CUBIC flow.
- (b) 1 BBRv2 flow, 10 CUBIC flows. (c) 10 BBRv2 flows, 10 CUBIC flows. (d) 50 BBRv2 flows, 50 CUBIC flows.

No losses

The Effects of Fixed-rate Pacing on TCP CUBIC

- 100 simultaneous flows
- Bottleneck is 1Gbps, 20ms propagation delay
- Pacing rate per flow is 8.5Mbps, 85% link utilization



Conclusions

- BBRv2 addressed the limitations of BBRv1
 - Tolerates much higher packet loss rates than CUBIC, lower than BBRv1
 - Mitigates the unfairness problem
 - Has better coexistence with CUBIC than BBRv1
 - Produces lower retransmission rates than BBRv1
 - > Exhibits low queueing delay, even with bloated buffers
 - Coexistence / fairness problem with CUBIC is not solved but mitigated

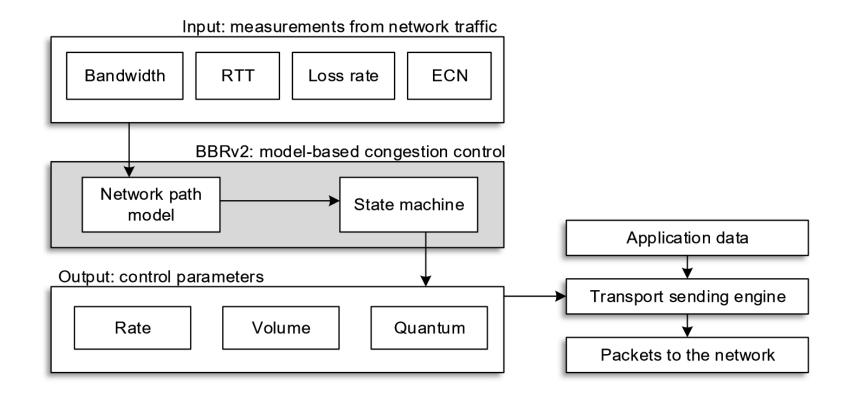
Resources

- Full paper: <u>https://www.sciencedirect.com/science/article/pii/S014036642030092X</u>
- The scripts for emulation are available in the following GitHub repository https://github.com/gomezgaona/bbr2
- BBRv2 version: v2alpha-2019-07-28
- Cyberinfrastructure website: <u>htttp://ce.sc.edu/cyberinfra/</u>
- BBRv2 GitHub repository: https://github.com/google/bbr/tree/v2alpha



BBRv2

BBRv2 architectural diagram¹

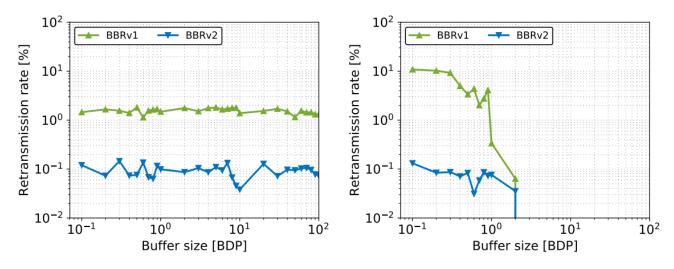


^{1.} N. Cardwell et al. "BBR v2, A Model-based Congestion Control." IETF 104, March 2019.



Retransmission Rate

- Two simultaneous flows, one with 10ms prop delay, and the other with 50ms
- Bottleneck is 1Gbps
- BBRv2 produces a low retransmission rate, (e.g., ~0.1%)

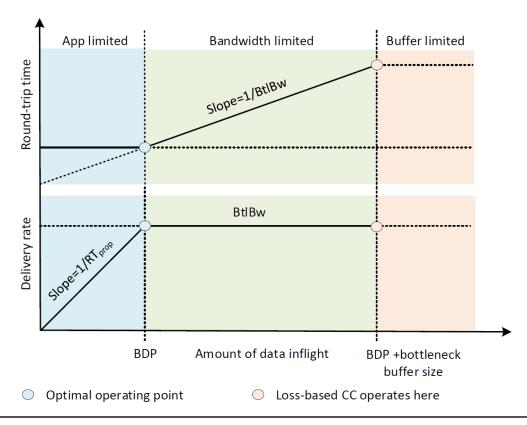


- (a) Retransmissions, Tail Drop.
- (b) Retransmissions, FQ-CoDel.

Retransmission rate generated by BBRv1 and BBRv2

BBR- Model-based Congestion Control

- BBR: Bottleneck Bandwidth and Round-trip propagation time¹
- Dynamically estimates windowed max BW and min RTT on each ACK



^{1.} N. Cardwell et al. "BBR: Congestion-Based Congestion Control: Measuring bottleneck bandwidth and round-trip propagation time." Queue 14.5 (2016): 20-53..